

CallExpr E

$T = \alpha$

Function: append

$T = ([\gamma], \gamma) \rightarrow [\gamma]$

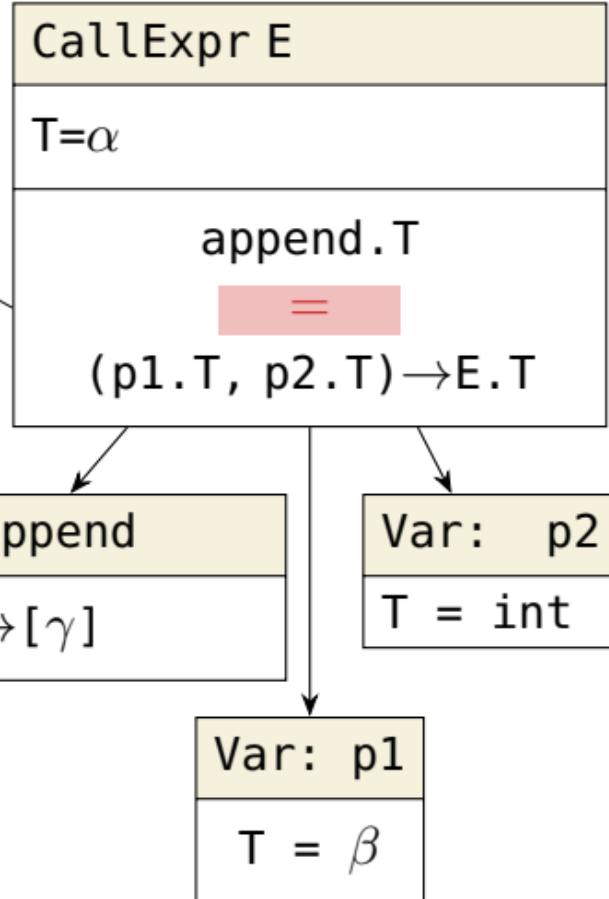
Var: p2

$T = \text{int}$

Var: p1

$T = \beta$

Semantische
Regel
für
CallExpr



CallExpr E

$T = \alpha$

$([\gamma], \gamma) \rightarrow [\gamma]$

=

$(\beta, \text{int}) \rightarrow \alpha$

Function: append

$T = ([\gamma], \gamma) \rightarrow [\gamma]$

Var: p2

$T = \text{int}$

Var: p1

$T = \beta$

Ersetzung der Typvariablen

CallExpr E

T=[γ]

([γ], γ) → [γ]

=

(β , int) → [γ]

Function: append

T=([γ], γ) → [γ]

Var: p2

T = int

Var: p1

T = β

Ersetzung der Typvariablen

$\alpha \mapsto [\gamma]$

CallExpr E

T=[γ]

([γ], γ) \rightarrow [γ]

=

([γ], int) \rightarrow [γ]

Function: append

T=([γ], γ) \rightarrow [γ]

Var: p2

T = int

Var: p1

T = [γ]

Ersetzung der Typvariablen

$\alpha \mapsto [\gamma]$

$\beta \mapsto [\gamma]$

CallExpr E

T=[int]

([int], int) → [int]

=

([int], int) → [int]

Function: append

T=([int],int)→[int]

Var: p2

T = int

Var: p1

T = [int]

Ersetzung der Typvariablen

$\alpha \mapsto [\gamma]$

$\beta \mapsto [\gamma]$

$\gamma \mapsto \text{int}$