

CallExpr E

$T = \alpha$

Function: append

$T = ([\gamma], \gamma) \rightarrow [\gamma]$

Var: p2

$T = \text{int}$

Var: p1

$T = \beta$

Semantische
Regel
für
CallExpr

CallExpr E

$T = \alpha$

append.T

=

$(p1.T, p2.T) \rightarrow E.T$

Function: append

$T = ([\gamma], \gamma) \rightarrow [\gamma]$

Var: p2

$T = \text{int}$

Var: p1

$T = \beta$

CallExpr E

$T = \alpha$

$([\gamma], \gamma) \rightarrow [\gamma]$

=

$(\beta, \text{int}) \rightarrow \alpha$

Function: append

$T = ([\gamma], \gamma) \rightarrow [\gamma]$

Var: p2

$T = \text{int}$

Var: p1

$T = \beta$

Ersetzung der Typvariablen

CallExpr E

$T=[\gamma]$

$([\gamma], \gamma) \rightarrow [\gamma]$

=

$(\beta, \text{int}) \rightarrow [\gamma]$

Function: append

$T=([\gamma], \gamma) \rightarrow [\gamma]$

Var: p2

$T = \text{int}$

Var: p1

$T = \beta$

Ersetzung der Typvariablen

$\alpha \mapsto [\gamma]$

CallExpr E

$T=[\gamma]$

$([\gamma], \gamma) \rightarrow [\gamma]$

=

$([\gamma], \text{int}) \rightarrow [\gamma]$

Function: append

$T=([\gamma], \gamma) \rightarrow [\gamma]$

Var: p2

$T = \text{int}$

Var: p1

$T = [\gamma]$

Ersetzung der Typvariablen

$\alpha \mapsto [\gamma]$

$\beta \mapsto [\gamma]$

CallExpr E

T=[int]

$([int], int) \rightarrow [int]$

=

$([int], int) \rightarrow [int]$

Function: append

T=([int],int)→[int]

Var: p2

T = int

Var: p1

T = [int]

Ersetzung der Typvariablen

$\alpha \mapsto [\gamma]$

$\beta \mapsto [\gamma]$

$\gamma \mapsto \text{int}$